18/11/2015 512/1/1

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محاضرة [7]

* Circular Convolution :-

If the sequences (Discrete time sequences) are periodic, then the Convolution is circular.

 $y(n) = X_1(n) N X_2(n)$ where $X_1(n) \notin X_2(n)$ are two periodic sequences every

N samples.

N > The symbol of circular convolution.

 $y(n) = X_1(n) \otimes X_2(n) = \sum_{m=0}^{N-1} X_1(m), X_2((n-m))_N$

 $=\sum_{n=0}^{N-1}x_{i}((n-m),x_{i}(m))$

= IDFT { X,(K), X2(K) }

DFT DFT

for X,(n) for X2(n)

Example: - Compute the circular Convolution for:

$$X_1(n) = \{2,1,2,1\} \text{ and } X_2(n) = \{1,2,3,4\}$$

for N=4 (Find X1(n) @ X2(n))

y(n)= X1(n) (4) x2(n) = 5 x1(m) x2(n-m))4

	M	2		1		
		0	(2	3	
	x, (m)	2	1	7		
	Xi(m)				,	
10 -		1	2	3	4	
11 = 0	×2 ((1-m))4	(41	3		
N = 1	X2 ((1-M))	2		3	2	⇒ 5
	1/2 (12-17)/4		\	4	3	-
		3	2	1	-	
11 = 3	X2 ((3-m))4	4	3		1 4	3
				2	1	(=)

11

$$y(0) = 2 \times 1 + (\times 4 + 2 \times 3 + (\times 2 = 14$$

 $y(1) = 2 \times 2 + (\times 1 + 2 \times 4 + 1 \times 3 = 16$
 $y(3) = 16$

- Another Method

For shifting
$$X_1(n)$$

 $y(n) = \sum_{m=0}^{\infty} X_2(m) X_1((n-m))_4$

$$X_{1}((-n))_{4} = 2 \cdot 1 \cdot 2$$

Another Solution: yon = x(n) (1) x(n)= IDFT (x(k), x, (k))

Dompute OFT for
$$X_{1}(n)$$
 & $X_{2}(n)$

$$X_{1}(k) = \begin{pmatrix} 6 \\ 2 \\ 3 \end{pmatrix}, X_{2}(k) = \begin{pmatrix} -2+2j \\ -2-2j \\ -2-2j \end{pmatrix}$$

$$X_{1}(k), X_{2}(k) = \begin{pmatrix} 6 \\ 6 \\ 7 \\ -2-2j \end{pmatrix} = \begin{pmatrix} 60 \\ -2+2j \\ -2-2j \end{pmatrix} = \begin{pmatrix} 60 \\ -2+2j \\ -2-2j \end{pmatrix}$$

2 Compat I'DFT y(n) - (14)

Prove the following properties:-

$$\Im W_N^2 = W_{N/2}$$

[2]
$$w_{N}^{K+N/2} = w_{N}^{K} (w_{N/2}^{N/2}) \rightarrow -1$$
, $w_{N}^{N/2} = e^{j} \pi$

$$\frac{131}{M} \frac{W_{N/2}}{W_{N/2}} = e^{-\frac{3211}{N_2}} = (e^{-\frac{3211}{N}})^2 = w_N^2$$

* Fast Fourier Trans form (FFT)

FFT Algorithms are many algorithms used to reduce the

Complex computation for DFT

U Radix - 2 DIT FFT

splitting Decimation in time

for periodic sequence x(n), the discrete fourier transform (DFT) is X(K), which has periodic N samples

$$X(K) = \sum_{n=0}^{N-1} x(n) W_N^{Kn} \Rightarrow \text{The previous method}$$
where $w_N = e^{\frac{i}{2}} \frac{2\pi}{N}$

For Radix- 2 DIT FFT algorithm:

$$X(R) = \sum_{n=\text{even}} + \sum_{p=\text{odd}} \frac{N}{2} \text{ point DFT}$$

$$X(|X|) = \sum_{n=0}^{N_2-1} X(2n) w_N$$

$$= \sum_{n=0}^{N_2-1} X(2n+1) w_N$$

$$= \sum_{n=0}^{N_2-1} X(2n+1) w_N$$

$$= \sum_{n=0}^{N_2-1} X(2n+1) w_N$$

$$= \sum_{n=0}^{N/2-1} \chi(2n) w_{N/2}^{kn} + w_{N}^{k} \sum_{n=0}^{N/2-1} \chi(n+1) w_{N/2}^{kn}$$

assume
$$f_1(n) = \chi(2n)$$

 $f_2(n) = \chi(2n+1)$

For Large values of N, this operation is repeated until we reach to 2-point DFT computation

Example:
$$X(n) = \{X(0), X(1)\} \implies N = 2$$

 $X(R) = \{X(n), X(1)\} \implies N = 2$
 $X(R) = \{X(n), X(1)\} \implies N = 2$

$$= \chi(0) (\omega_{2}^{0}) + \chi(1) \omega_{2}^{K}$$

$$K=0 \Rightarrow \chi(K) = \chi(0) + \chi(1)$$

In general:

A =
$$a + b w_{p}^{r}$$

b

Wr

B = $a - b w_{p}^{r}$

General form for butterfy computation

Ex: FFT $\frac{f_{or} N=4}{X(K)=F_{1}(K)+W_{N}^{K}F_{2}(K)} \Rightarrow 0$

 $X(K+\frac{N}{2}) = F_1(K) - w_N^K F_2(K) \Rightarrow 0$ $K=0, 1, -, \frac{N}{2}-1$

where $f_1(n) = \chi(2n) \rightarrow \text{even numbered seq.}$ $f_2(n) = \chi(2n+1) \rightarrow \text{odd}$

 $X(K) = F_1(K) + \omega_4^K F_2(K) \longrightarrow 0$ $X(K+2) = F_1(K) - \omega_4^K F_2(K) \longrightarrow 0$ K = 0, 1

- Turn over

ear (1) K=0 => X(0) = F1(0) + Wy F2(0) ? 2-point DFT $K=1 \Rightarrow X(1) = F_1(1) + \omega_4 F_2(1)$ ey(2) K=0 => X(2) = F,(0) - W4 F2(0) & 2-point DFT K=1 => X(3) = F,(1) - w4 F2(1) (0) X(0) Whit DF X(0) = f(0) [0] X(2)= f(1) [3] 1-2+2j/ X(1) (-2) F1(1) X(1)= f2(4) (2) X (2) X(3)= f2(1) [3] $f_{1}(n) = \chi(2n) \Longrightarrow n = 0.1$ $f_{i}(n) = \{ \chi(0), \chi(z) \}$ $f_{z}(n) = x(zn+1) \Rightarrow n=0,1$ coldfor (n) = { x(1), x(3) } exi find Radix - 2 DIT FFT for X(n)= { 0,1, 2,3 } X(K) $\{6, -2+j2, -2, -2-2j\}$ اكل مسقَّط على الرحمة الله بنم باللوم الأم